

A Variable Neighborhood Search approach for the S-labeling problem



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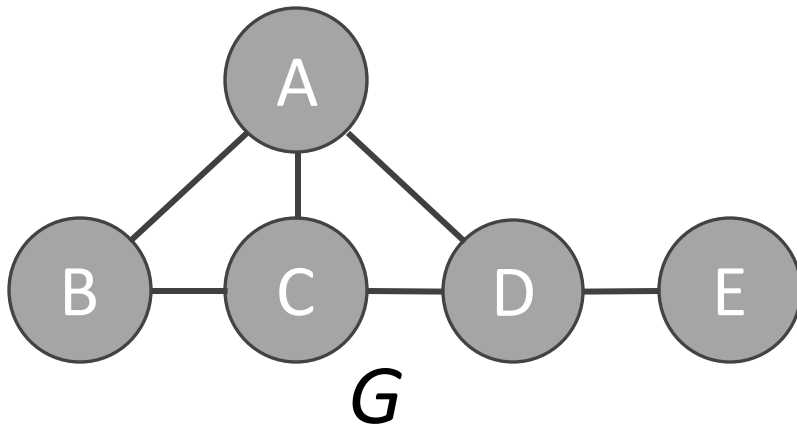
1. Introduction
2. The S-labeling problem
3. Previous works
4. Our proposal
5. Algorithmic results
6. Conclusions and future work

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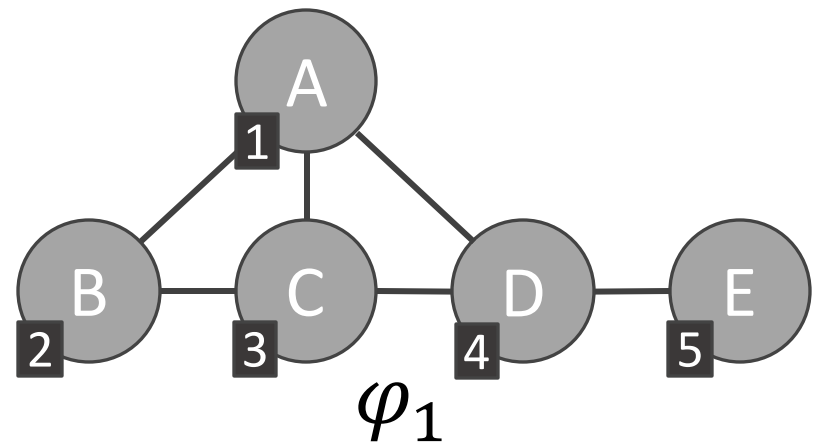
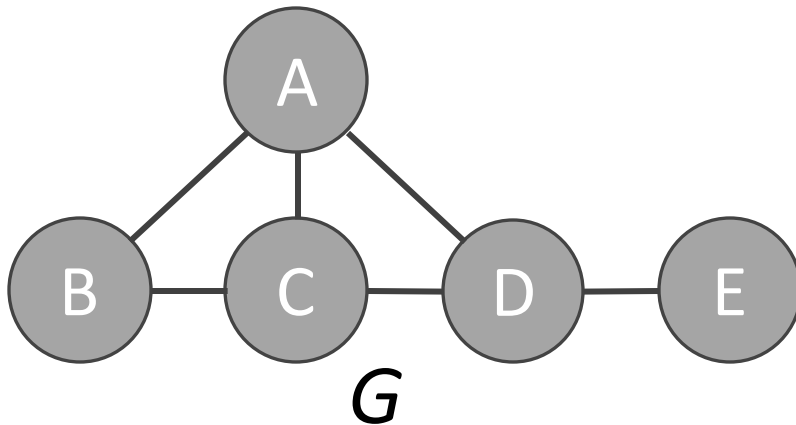
Introduction – Graph Labeling Problems

- Graph Labeling Problems are a kind of **combinatorial optimization problems**.
- A labeling of a graph consist of assigning **labels** to each vertex of an input graph to **optimize** a certain objective function.



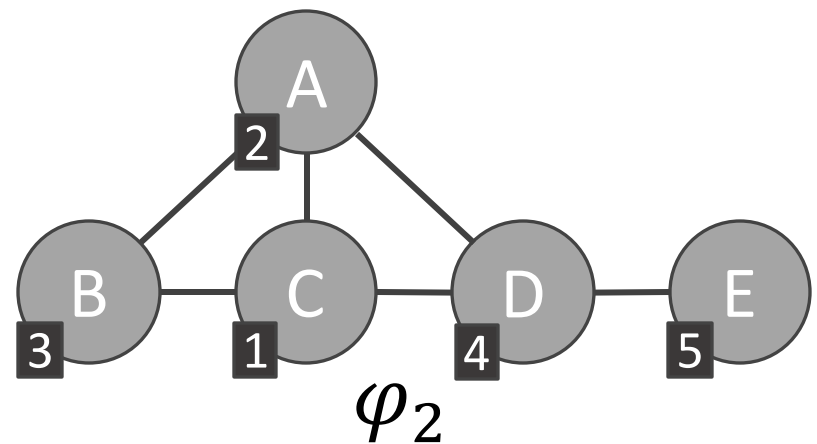
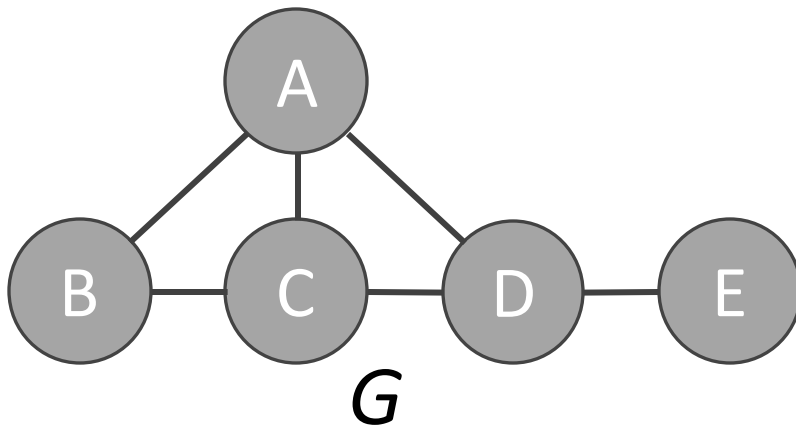
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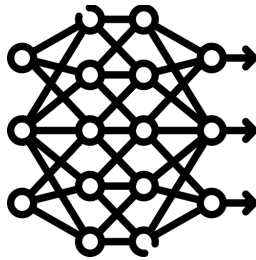
Introduction – Graph Labeling Problems

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Introduction – Graph Labeling (2)

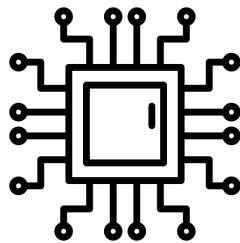
- **Graph Labeling Problems** has been found to have a lot of **real-world applications**:



Network optimization



Numerical analysis



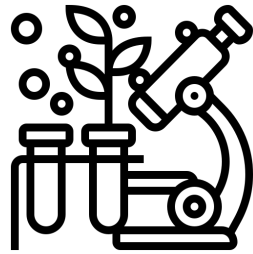
Circuit desing



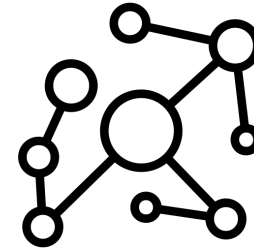
Information retrieval

Introduction – Graph Labeling (3)

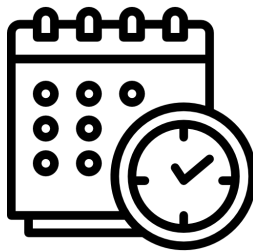
- **Graph Labeling Problems** has been found to have a lot of **real-world applications**:



Computational biology



Graph theory



Scheduling



Archaeology

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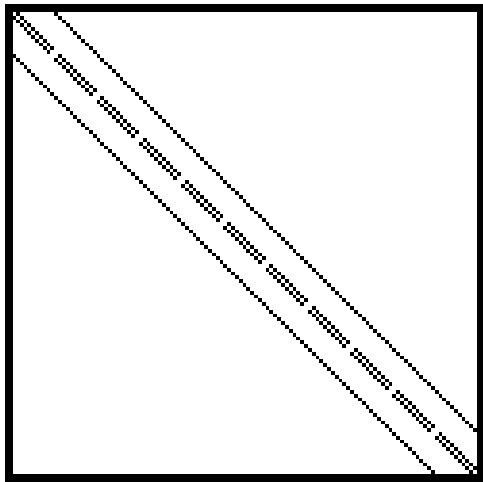
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Practical applications for the S-labeling

- The S-labeling problem was originally proposed in the context of the matrix packaging [7].
 - Matrix packaging consists of **permutating** the rows and columns of a sparse matrix to make **calculations or storage easier**.
- The S-labeling problem is a specific case of matrix packaging in which the matrix is **zero trace symmetric (0, 1)-matrix**, which represents an undirected graph.

Practical applications for the S-labeling

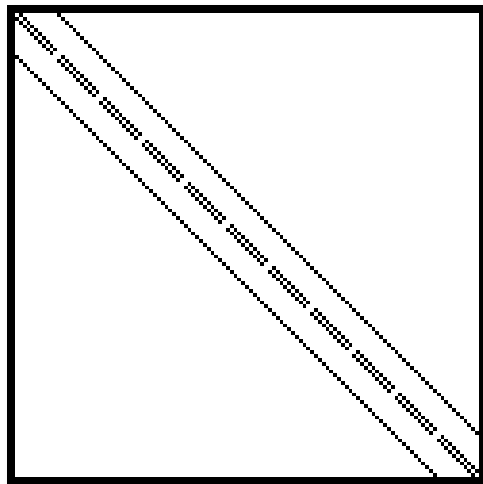
- By applying the S-labeling to a sparse matrix we get a new one that is **easier to compute**.
- **Example** of one of the instances used:



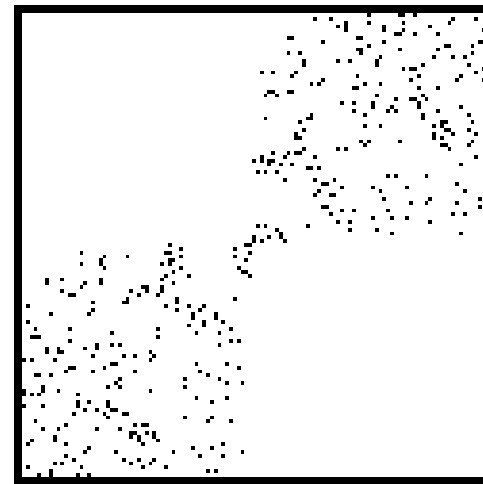
Before

Practical applications for the S-labeling

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- **Example** of one of the instances used:



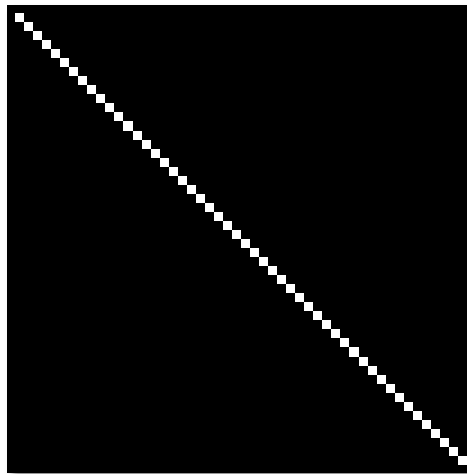
Before



After

Practical applications for the S-labeling (2)

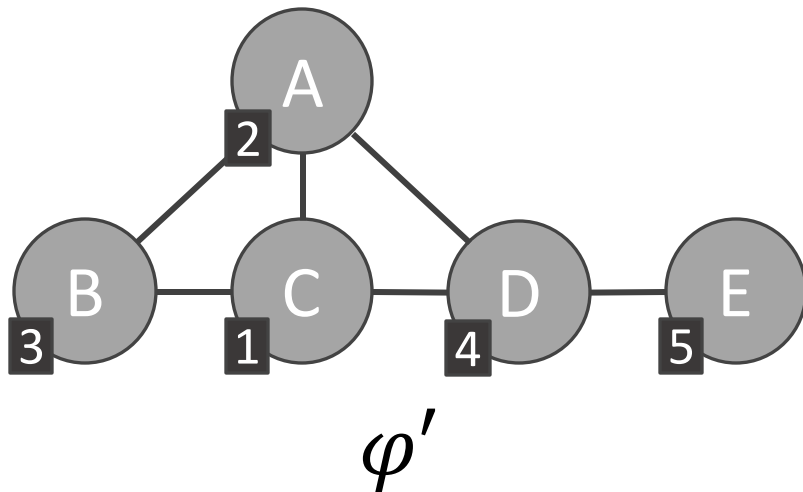
- The S-labeling is **only useful** when applied to **sparse** matrices.
- For example, the solution for **complete graphs** is **trivial**.



Complete graph

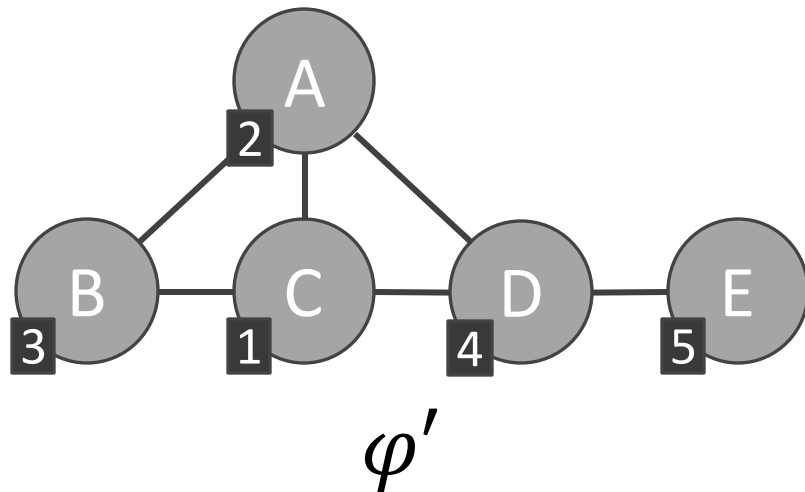
Problem description

- Given a graph labeling φ , we define the **objective function value** as the sum of the evaluation of each edge.
- We evaluate an edge as the **minimum label** assigned to the vertices of that edge.



Problem description

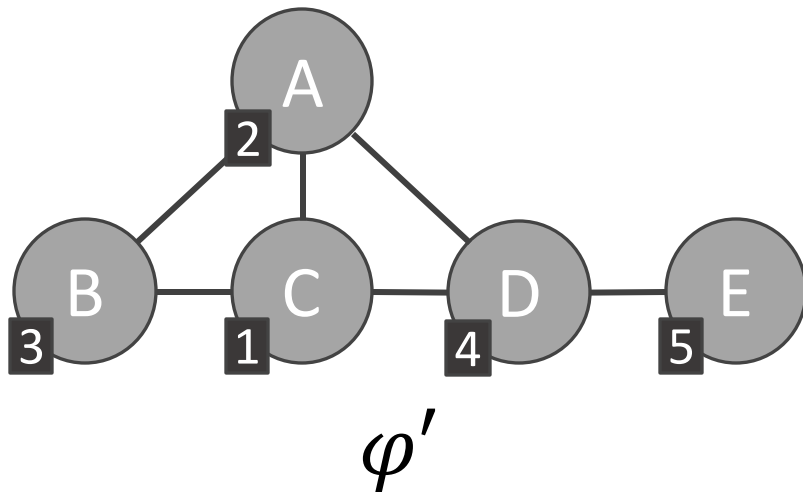
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$$Eval(\varphi', (A, B)) = \min(2, 3) = 2$$

$$Eval(\varphi', (A, C)) = \min(2, 1) = 1$$

$$Eval(\varphi', (A, D)) = \min(2, 4) = 2$$

$$Eval(\varphi', (B, C)) = \min(3, 1) = 1$$

$$Eval(\varphi', (C, D)) = \min(1, 4) = 1$$

$$Eval(\varphi', (D, E)) = \min(4, 5) = 4$$

Problem description (2)

- Given a graph labeling φ , we define the **objective function value** as the sum of the evaluation of each edge.
- We evaluate an edge as the **minimum label** assigned to the vertices of that edge.

$$\left. \begin{array}{l} Eval(\varphi', (A, B)) = \min(2, 3) = 2 \\ Eval(\varphi', (A, C)) = \min(2, 1) = 1 \\ Eval(\varphi', (A, D)) = \min(2, 4) = 2 \\ Eval(\varphi', (B, C)) = \min(3, 1) = 1 \\ Eval(\varphi', (C, D)) = \min(1, 4) = 1 \\ Eval(\varphi', (D, E)) = \min(4, 5) = 4 \end{array} \right\} \sum_{(u,v) \in E} Eval(\varphi', (u, v)) = \mathbf{11}$$

Problem description (3)

- Given a graph labeling φ , we define the **objective function value** as the sum of the minimum label of the vertices of each edge.

$$O.F.(\varphi') = \sum_{(u,v) \in E} \min(\text{label}(\varphi', u), \text{label}(\varphi', v))$$

- The objective in the S-labeling is to find the labeling φ^* among all the labelings ϕ that **minimizes the objective function**.

$$\varphi^* = \arg \min_{\varphi \in \Phi} O.F.(\varphi)$$

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Previous works

2006

- The problem is first proposed. [7]

2018

- The problem is studied **theoretically**. [2]

2019

- An **exact solution method** based on MIP is proposed. [6]

2023

- A heuristic approach based on a **Population-based Iterated Greedy** is presented. [5]

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Why VNS?

- Multiple **population-based methods** have already been studied.
 - We want to study other methods other than population-based metaheuristics.
- VNS have **multiple variants**, that fit different situations.
- We have **already used VNS** methods on other problems successfully.

Summary of our proposal

1 Random constructive method.

3 Different Shake methods

- Shuffle, random movement, and inverse.

2 Local Searches

- Swap First Improvement and Insert First-Best.

3 VNS variants

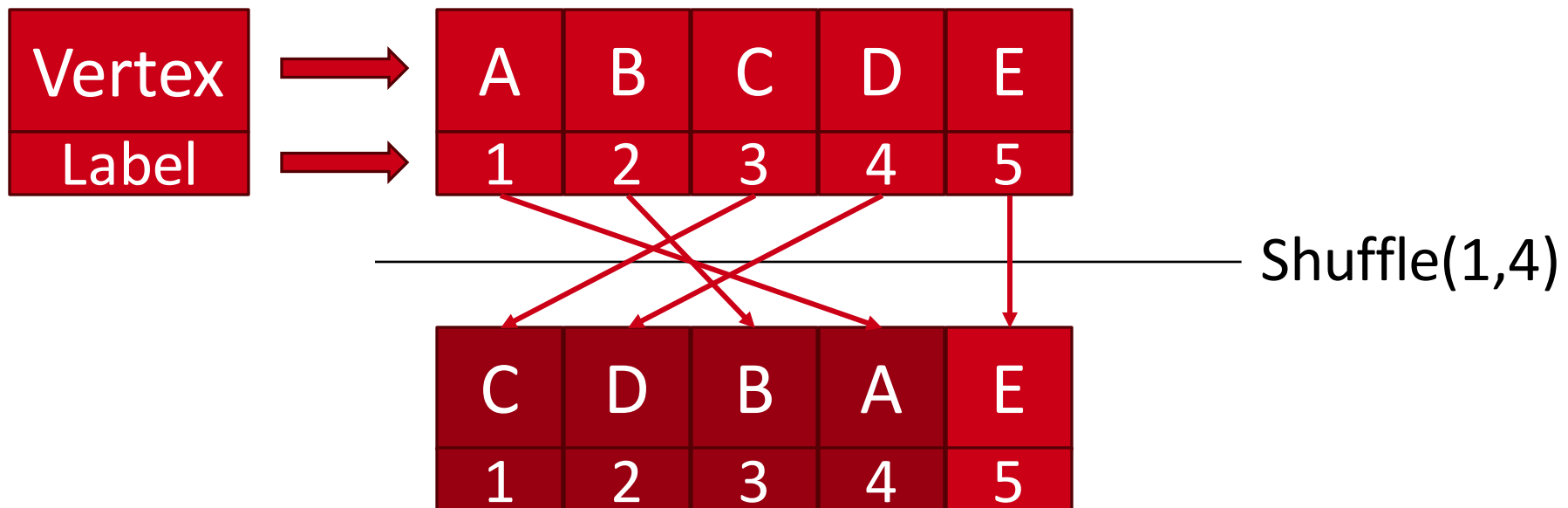
- BVNS, VND and GVNS.

VNS variants

- **Basic VNS (BVNS)**: applies the VNS schema without any modification.
- **Variable Neighborhood Descent (VND)**: removes the Shake step and adds another Local Search to scape from local optimums.
- **General VNS (GVNS)**: replaces the Local Search step in BVNS for a VND.

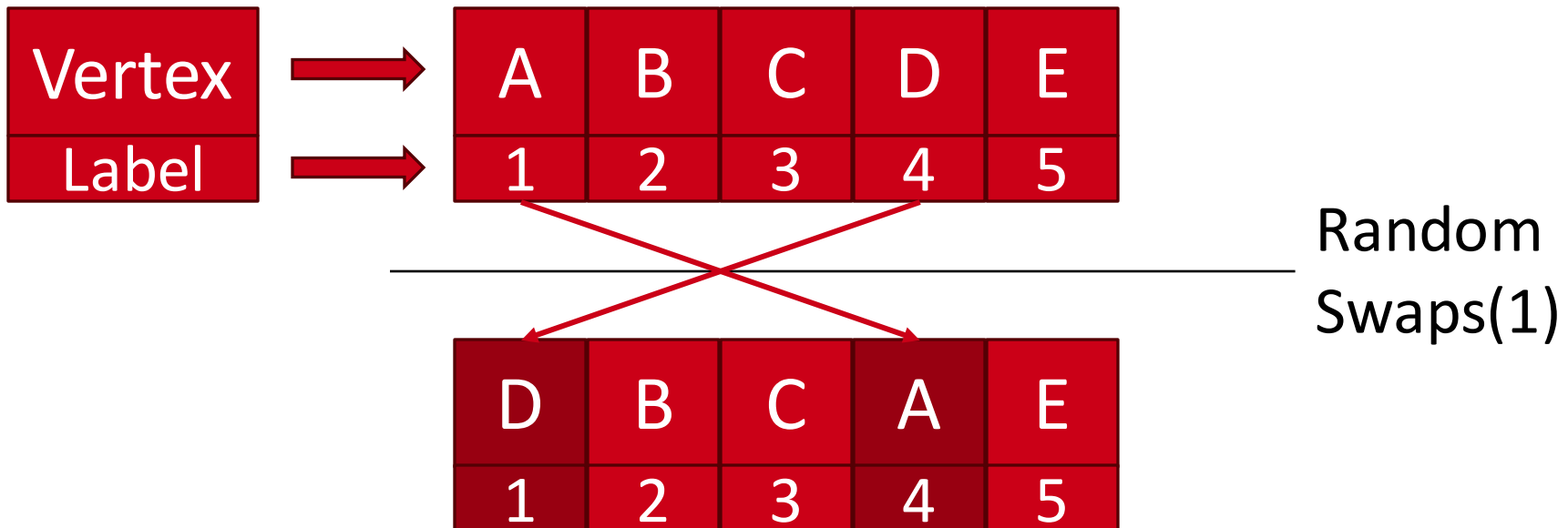
Shake methods (1)

- **ShuffleShake**: randomly shuffling the label of a subset of vertices with labels in the range $[1, k]$.



Shake methods (2)

- **NeighborhoodShake**: execute k random swap and insert movements.



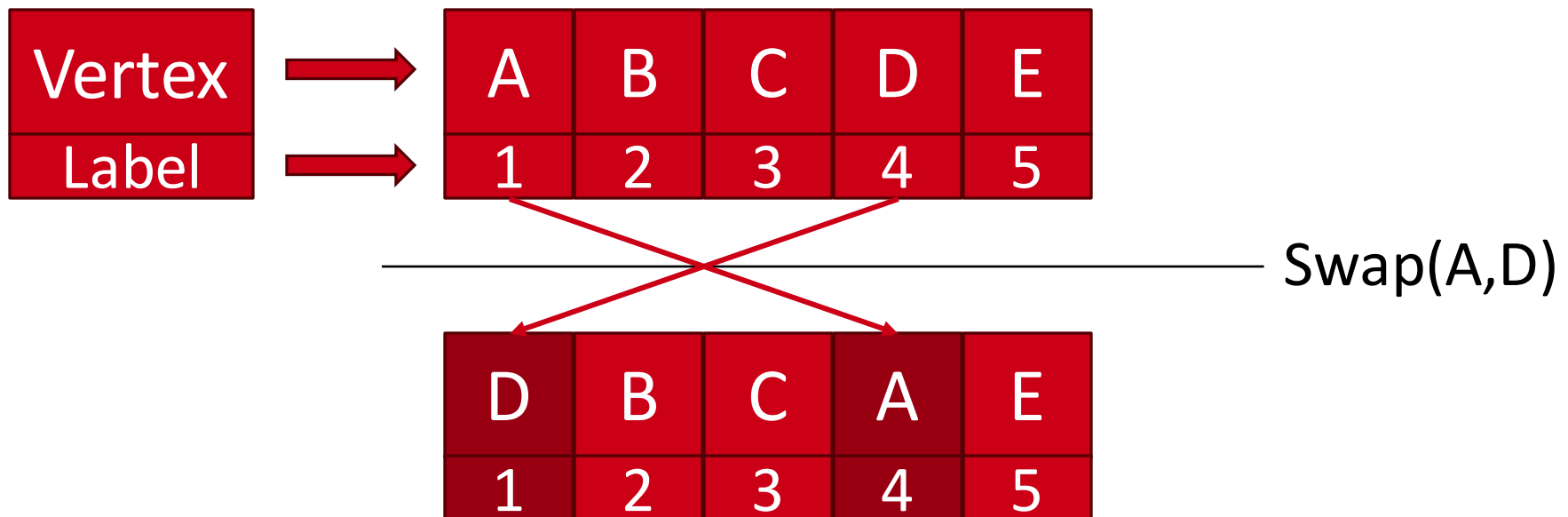
Shake methods (3)

- **InverseShake**: assign the highest labels to the vertices with the lowest initial labels, and vice versa.



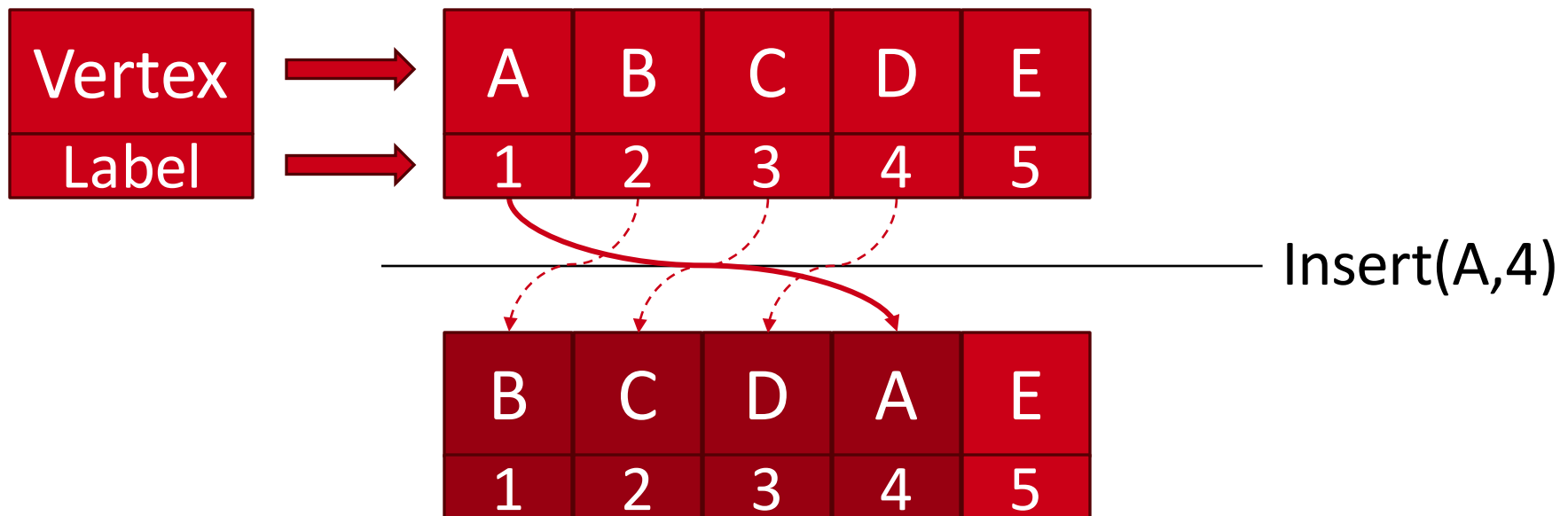
Local Search methods – Movements (1)

- **Swap movement:** exchange the assigned labels of two vertices.



Local Search methods – Movements (2)

- **Insert movement:** assign a certain label to a given vertex, displacing the rest of vertices.



Local Search methods - Strategy

- The most common ones are
 - **First Improvement (FI)**: commit the first movement that produces an improvement.
 - **Best Improvement (BI)**: commit the best movement among all possibles.
- For both movements the **BI** strategy was tested and found **too costly**.
- For the **Swap** movement we chose the **FI strategy**, which produced good results and diversified the search.

Local Search methods – Insert strategy

- For the insert movement we also propose a **First-Best Strategy** for the **Insert movement**.
- This new strategy surged from the idea that the insert movement produces **multiple intermediate states** which **can be evaluated**.



Summary of our proposal

1 Random constructive method.

2 Local Searches

- Swap First Improvement and Insert First-Best.

3 Different Shake methods

- Shuffle, random movement, and inverse.

3 VNS variants

- BVNS, VND and GVNS.

• **A total of 14 different combinations.**

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 2. Comparison with the state-of-the-art method
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BVNS – Tuning

Shake	Inverse	Movement	Shuffle	Inverse	Movement	Shuffle
LS	Insert			Swap		
Avg. O.F.	1539906	1687174	1548403	1489568	1576198	1489481
CPU T. (s)	300	300	300	300	300	300
% Dev.	1.77	10.58	1.95	0.04	4.48	0.04
% Best	0	0	0	45	0	55

- The most effective shake is the **Shuffle**.
- The most effective movement is the **Swap**.
- The best variant is **Shuffle + Swap**.

VND – Tuning

	VND – Insert&Swap	VND – Swap&Insert
Avg. O.F.	1606570	1629959
CPU T. (s)	303	301
% Dev.	0	1.15
% Best	100	0

- The most effective VND is the **Insert&Swap**.

GVNS – Tuning

Shake	Inverse	Movement	Shuffle	Inverse	Movement	Shuffle
VND	VND – Insert&Swap			VND – Swap&Insert		
Avg. O.F.	1505096	1524217	1496944	1508176	1531479	1488297
CPU T. (s)	300	300	300	300	300	300
% Dev.	0.67	1.57	0.25	0.74	1.93	0.02
% Best	10	0	15	0	0	75

- The most effective shake is the **Shuffle**.
- The most effective VND is the **Swap&Insert**.
- The best variant is **Shuffle + Swap&Insert**.

Comparison with the state-of-the-art

	State-of-the-art	BVNS	VND	GVNS
	Population-based Iterated Greedy (PIG)	Shuffle + Swap	Insert&Swap	Shuffle + Swap&Insert
Avg. O.F.	1477107	1489481	1606570	1488297
CPU T. (s)	200	300	303	300
% Dev.	0.00	0.72	6.17	0.60
% Best	95	5	0	0

- Among the proposed **VNS methods**, the best one is the GVNS and the worst one the VND.
- The VNS results are **close** to those of the PIG method, but they are still **outperformed by the state-of-the-art**.

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Conclusions

- We presented 14 different combinations of **VNS method** for the **S-labeling problem**. The **state-of-the-art** algorithm obtains **better results** than our proposal.
- The **GVNS** emerged as the **most effective**, with a **deviation lower than 1%**.
- The **VND** emerged as the **least effective**, showing that in this problem using **effective shake** methods is more important than using more local searches.

Future work

- Explore **new neighborhoods**, such as ejection chain.
- Explore alternative **constructive methods**.
- Implement methods which make use of the S-labeling **theoretical properties**.

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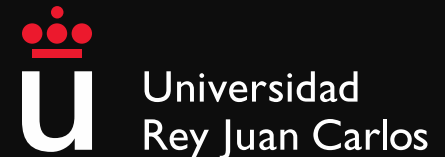


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